

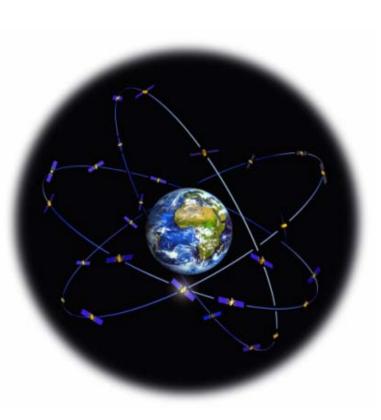
Decoding Galileo and Compass



Grace Xingxin Gao The GPS Lab, Stanford University June 14, 2007



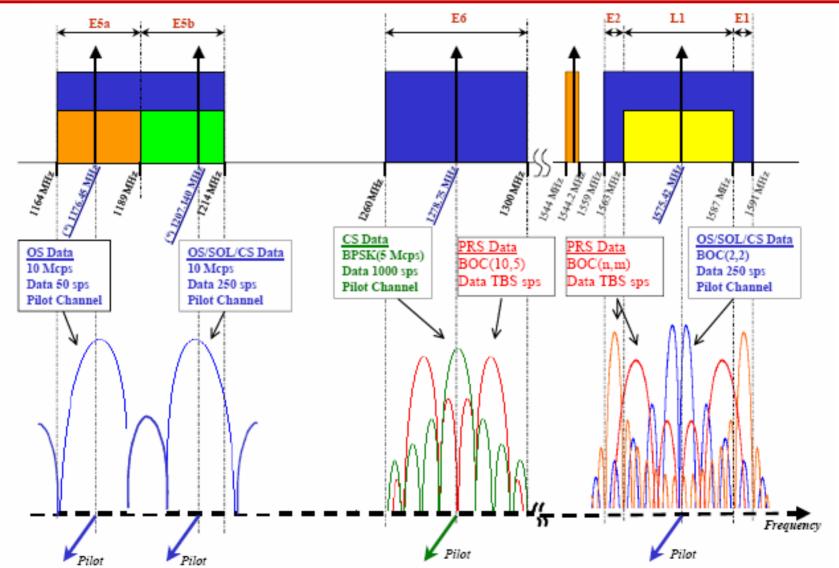
What is Galileo System?



- Global Navigation Satellite System
 built by European Union
 - The first Galileo test satellite GIOVE-A was launched on Dec. 28, 2005
 - First navigation signals were transmitted by GIOVE-A on Jan. 12, 2006
- Interoperable with GPS
- 30 satellites in three Medium Earth Orbit MEO planes at 23,222km above the earth
 - 9 satellite + 1 spare per plane
 - The inclination of the orbits was chosen to ensure good coverage of polar latitudes
 - One revolution 14 hours 4 min



Galileo Signal Spectrum





Why Seek the Galileo Codes?

- Galileo provides additional 30 satellites to the US GPS system
 - More accuracy and integrity
- Galileo L1 band overlaps with GPS L1 band
 - Can use the same antenna for the integrated Galileo/GPS receivers
- A Fancy Galileo Receiver



We are considering designs for new
signals from GPS and terrestrial
ranging sources that could augment
GPS. Thus, we are eager to gain a
deep understanding of the recent
efforts by our European colleagues.



Data Collection – My Antennas





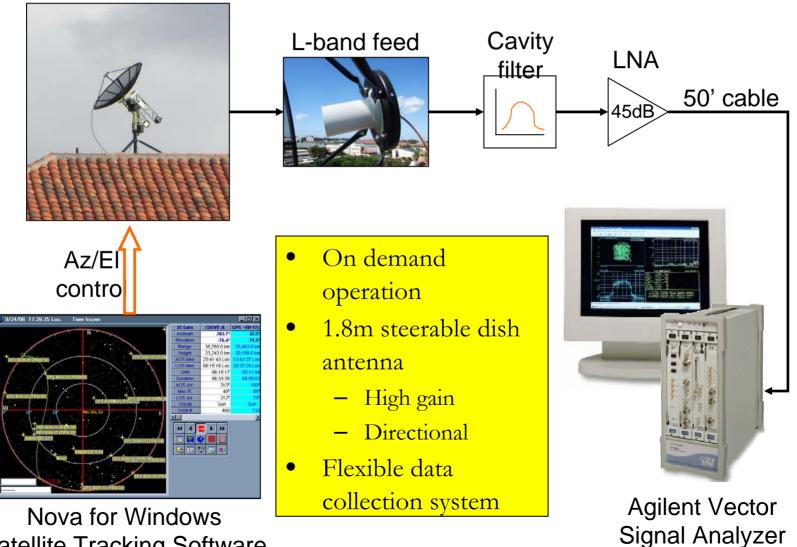
Data Collection - Antennas





Satellite Tracking Software

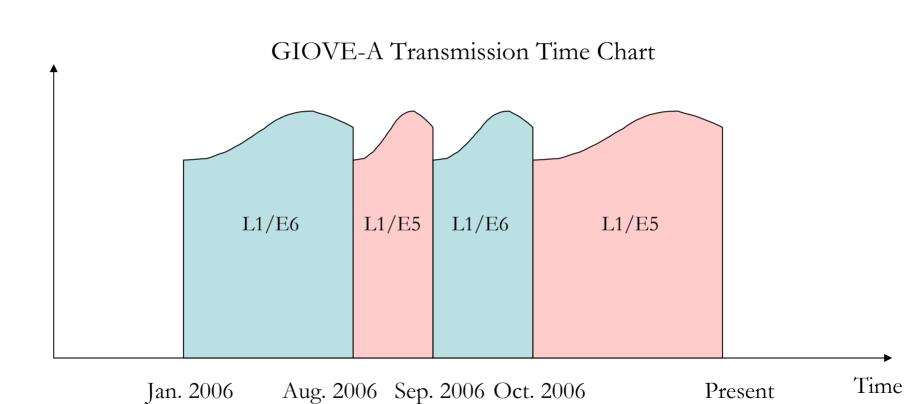
Data Collection Stanford GNSS Monitor Station



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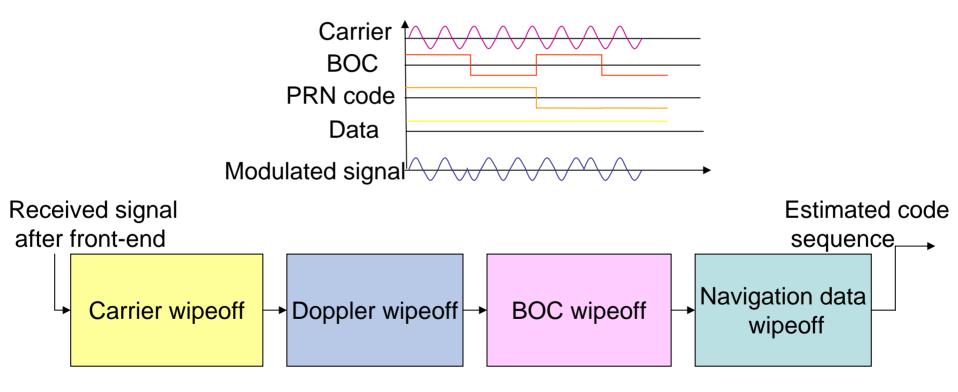
Seek L1 and E6 Codes





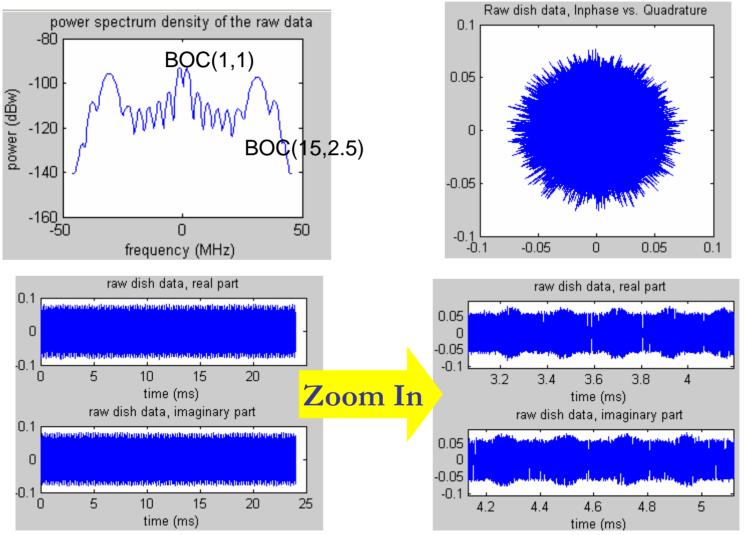
Estimate Individual Code Sequences

• Modulated signal is the product of carrier, BOC code, PRN code, and data



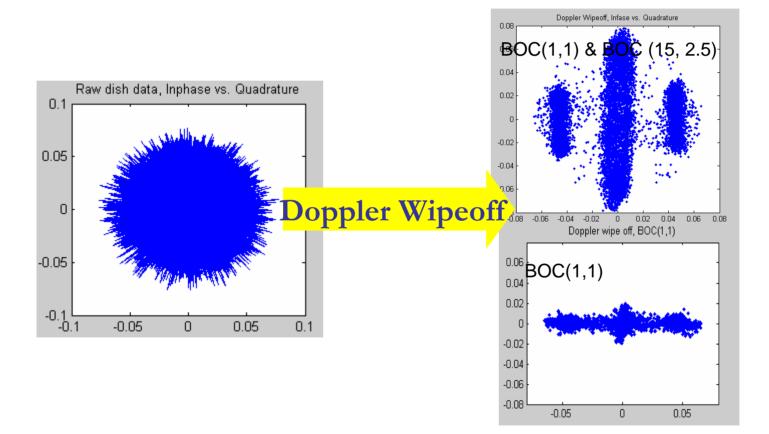


Raw Data Received from the Stanford Big Dish



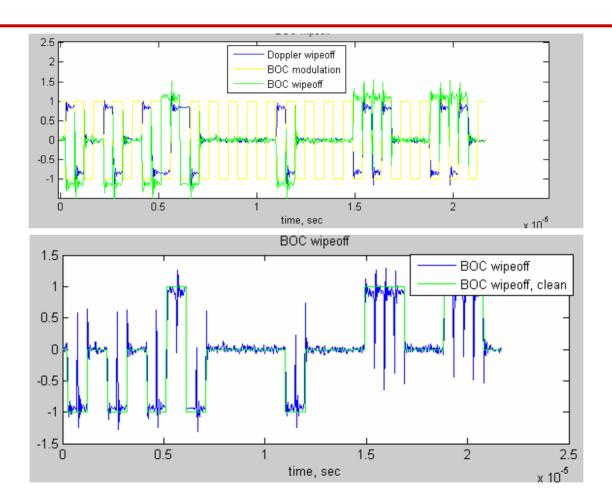


Doppler Wipeoff & Carrier Phase Alignment





BOC Wipeoff

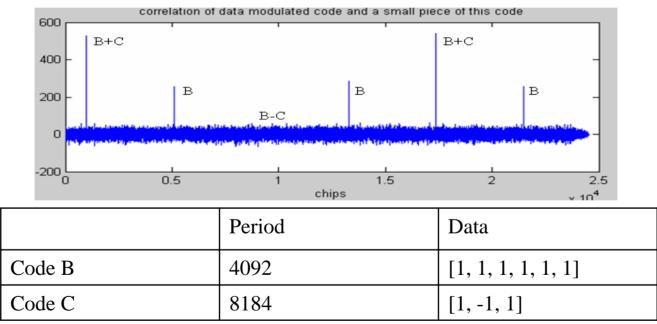


Align BOC code with the Doppler removed signal, and multiply them to wipe of the BOC code 12



Navigation Data Wipeoff

- Correlate the code sequence with a small slice of it
- The correlation plot shows that there are two codes, code B and code C with different periods.

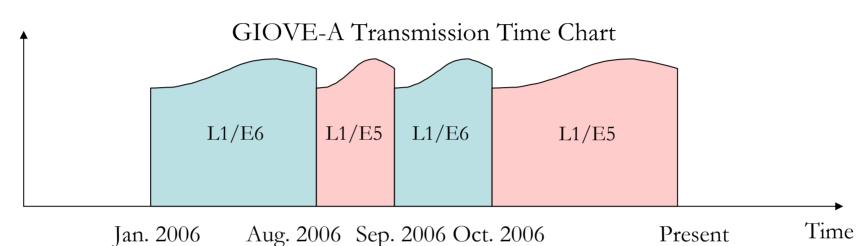


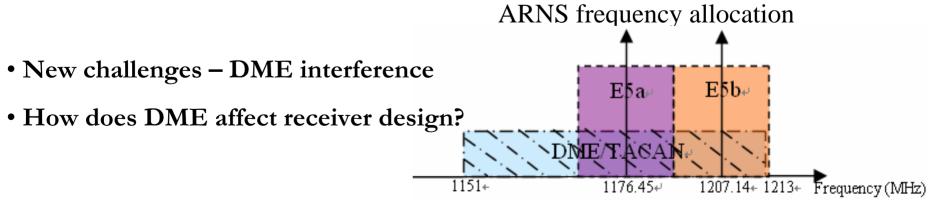
Code sequences are obtained after wiping off Doppler, BOC and navigation data

 Note, the sign of code B or code C is possibly flipped. This can be verified after the PRN code polynomials are calculated. If the polynomials contain "1+X", then the code is flipped. Otherwise, it's not.
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Seek E5a/E5b Codes







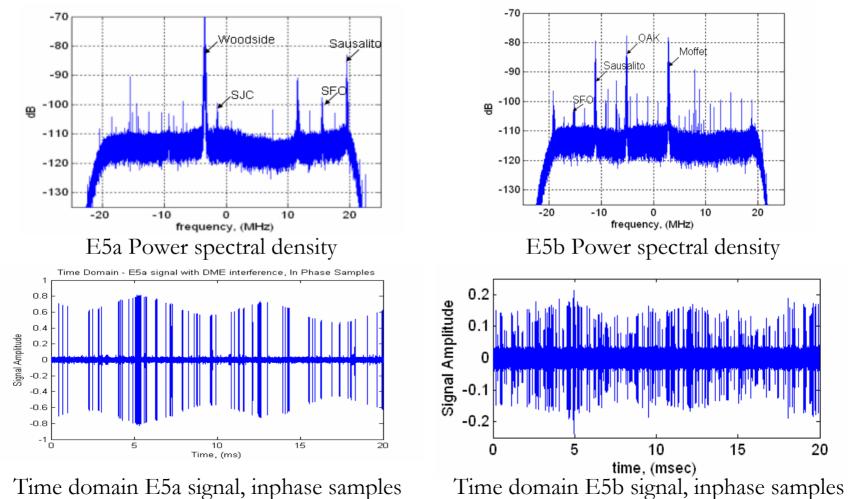
E5 Signal Environment

E5 signal environment in the bay area – DME/TACAN inference from nearby airports



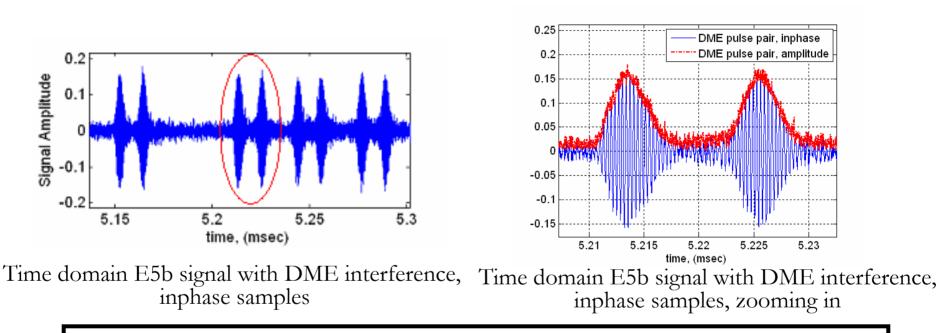


Received Raw Data





Raw Data, Zoom In

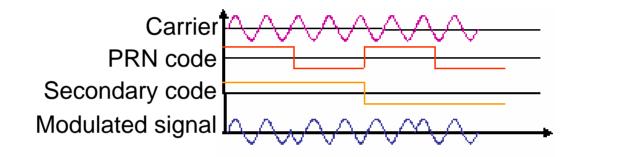


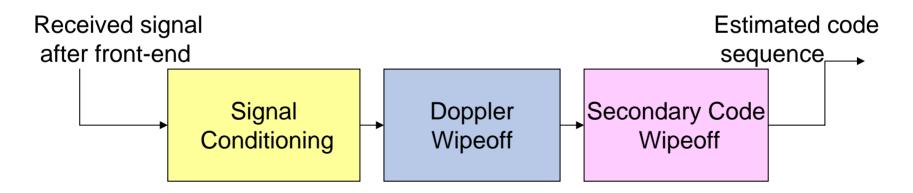
- The DME pulse amplitude is 5~100 times greater than the noise floor
- E5 signal is even buried in noise
- DME interference occurs 10-14% of the time
- \bullet DME comes in pairs with an inter-pulse interval of 12 μsec



Estimate Individual Code Sequences

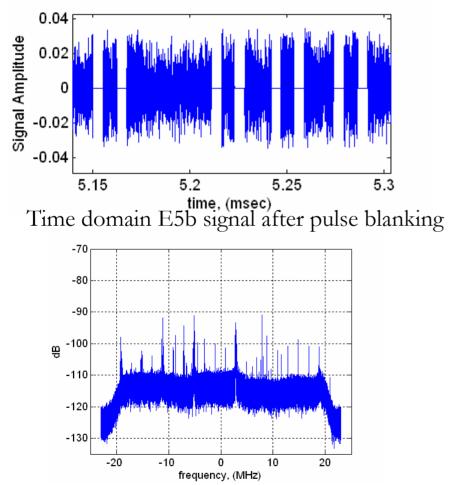
• Modulated signal is the product of a carrier, a PRN code, and a secondary code



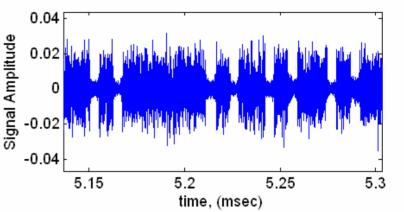




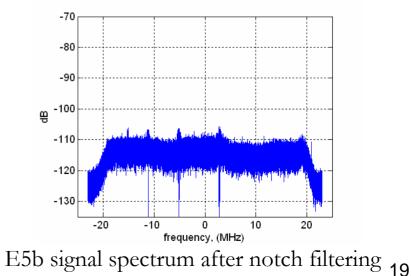
Signal Conditioning, **Pulse Blanking & Notch Filtering**



E5b signal spectrum after pulse blanking

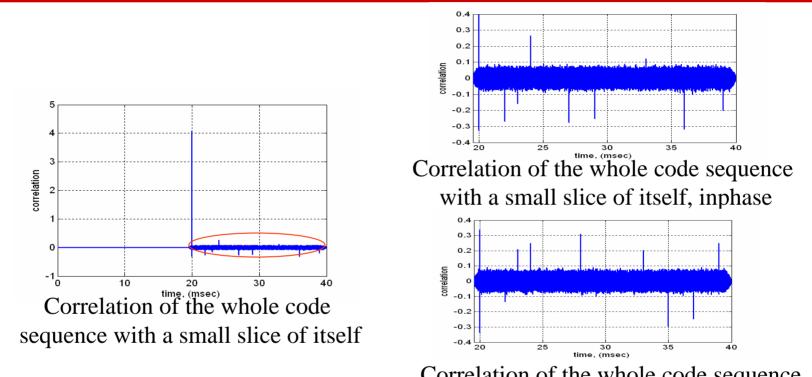


Time domain E5b signal after notch filtering





Calculate Code Periods

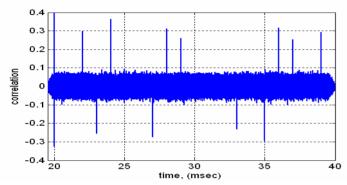


Correlation of the whole code sequence with a small slice of itself, quadrature

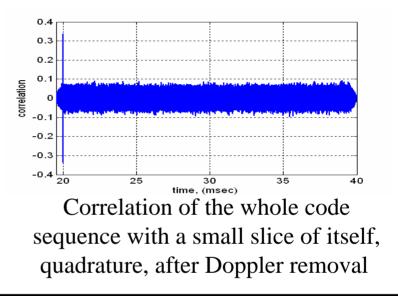
- Correlate the whole code sequence with a small slice of itself
- The intervals between pairs of peaks indicate a code period of 1ms
- Doppler results in constant phase variation, creating peaks in the inphase and quadrature channels and causing the peak heights to vary.



Doppler Wipeoff



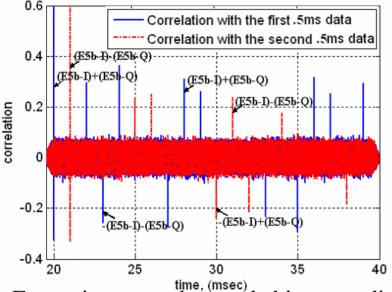
Correlation of the whole code sequence with a small slice of itself, inphase, after Doppler removal



- We search the whole Doppler domain from -5000 Hz to 5000 Hz and minimize the peak height variation after Doppler compensation
- After Doppler wipeoff and initial-phase adjustment, peaks with more uniform heights appear in the inphase channel and no peak in the quadrature channel.
- Null peaks indicate two codes added on top of each other



Secondary Code Wipeoff



	Period	Secondary code
E5b-I code	10230	[1, 1, 1, -1, 1, 1, 1, -1, 1, 1, -1, 1, -1, -1, 1, -1, 1, 1, -1, 1]
E5b-Q code	10230	$\begin{bmatrix} 1, -1, 1, -1, 1, -1, -1, -1, 1, 1, \\ 1, -1, 1, -1, -1, -1, 1, 1, 1, 1, 1, \end{bmatrix}$

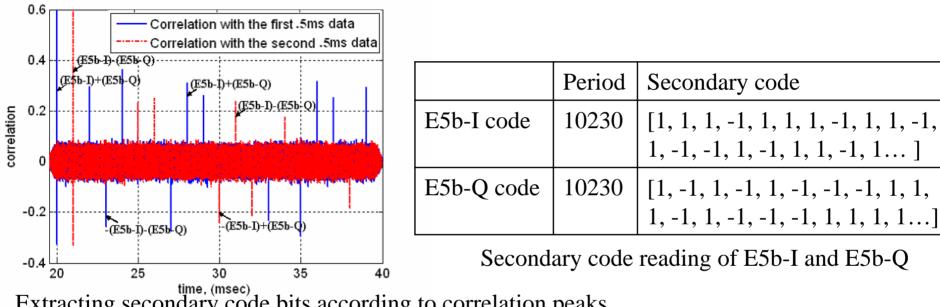
Secondary code reading of E5b-I and E5b-Q

Extracting secondary code bits according to correlation peaks

- Choose the slice length to be 0.5 msec to avoid bit transition within the small slice of data. Compute two consecutive slices of 0.5 msec and use the one with relatively larger peaks.
- Pick another slice of data that corresponds to a null peak to solve the ambiguity
- The sign of E5b-I code or E5b-Q code is possibly flipped. This can be verified after the PRN code polynomials are calculated. If the polynomials contain "1+X", then the code is flipped. Otherwise, it's not.



Secondary Code Wipeoff



Extracting secondary code bits according to correlation peaks

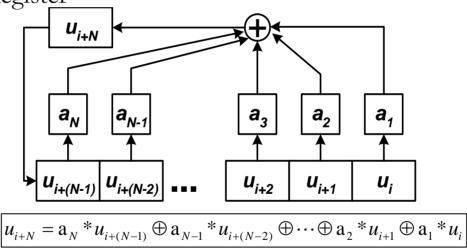
• With secondary codes available, we can coherently add multiple periods of signal together until the SINR ratio is high enough to boost the code chips above the noise floor

PRN code sequences are obtained!



Calculate Code Generator Stages (1/2)

- Start with linear codes
 - Linear codes can be generated by linear shift registers (LSR)
- Linear Shift Register

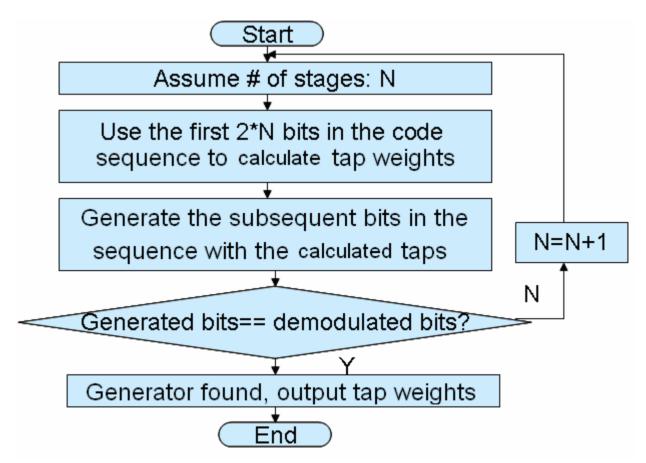


- 2*N bits contain all information of the whole 2^N-1 bits sequence
 - We have ui, i = 1, ..., 2*N, ...
 - The relationship for N distinct values of i yields N equations
 - Tapping weights are N unknown variables in N equations
 - Solve the equation, determine which taps are on or off (an = 0 or 1)



Calculate Code Generator Stages (2/2)

Calculate generator tap weights





Obtain Code Generators, L1 Signal

- Obtain code generator polynomials
 - Both codes are linear codes
 - Each code is generated by a 26 order polynomial
 - The obtained code generator polynomials can be factorized
 - Ambiguity of sign flipping is solved
 - The code sequence can be generated by modulo 2 adding maximal length sequences of the factor polynomials **Gold codes!**
- Code B (Gold code)
 - Poly1_code1 = $X^{13}+X^{10}+X^9+X^7+X^5+X^4+1$, initial state 111111111111
 - Poly2_code1 = $X^{13}+X^{12}+X^8+X^7+X^6+X^5+1$, initial state 1101110000011
- Code C (Gold code)
 - Poly1_code2 = $X^{13}+X^{10}+X^9+X^7+X^5+X^4+1$, initial state 1100110000011
 - Poly2_code2 = $X^{13}+X^4+X^3+X+1$, initial state 111111111111
- Secondary code of C
 - 25 bits long



Obtain Code Generators, E6 Signal

- E6-B code: 5115 bits, 1ms, gold code
 - Poly_E6_B_1: X^13+X^12+X^11+X^1+1;
 - initial state: 0 1 0 1 0 1 1 1 0 0 0 0 0
 - Poly_E6_B_2: $X^13+X^10+X^8+X^5+1$;
 - initial state: 1 1 1 1 1 1 1 1 1 1 1 1 1 1
- E6-C code: primary code* secondary code=511500 bits, 100ms Primary code: 10230 bits, 2ms, gold code
 - $Poly_E6_C_1: X^{14} + X^{8} + X^{7} + X^{4} + X^{3} + X^{2} + 1;$
 - initial state: 0 1 1 0 1 0 0 0 0 1 1 1 0 1
 - Poly_E6_C_2: X^14+X^11+X^6+X^1+1;
 - initial state: 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
 - Secondary code: 50 bits



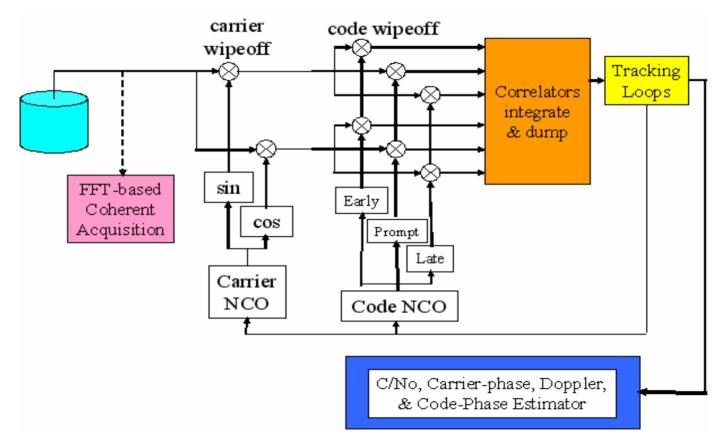
Obtain Code Generators, E5 signal

E5b-I code (10230 bits, 1msec, 14-stage Gold code)		E5a-I code (10230 bits, 1msec, 14-stage Gold code)	
Polynomial_1	X ¹⁴ +X ¹³ +X ¹¹ +X ⁴ +1	Polynomial_1	X ¹⁴ +X ⁸ +X ⁶ +X+1
Initial State_1	[1 1 1 1 1 1 1 1 1 1 1 1 1]	Initial State_1	[1 1 1 1 1 1 1 1 1 1 1 1 1]
Polynomial_2	X ¹⁴ +X ¹² +X ⁹ +X ⁸ +X ⁵ +X ² +1	Polynomial_2	X ¹⁴ +X ¹² +X ⁸ +X ⁷ +X ⁵ +X ⁴ +1
Initial State_2	[1 1 1 0 0 0 1 0 1 0 0 0 1 0]	Initial State_2	[1 1 1 0 1 0 1 0 1 1 1 1 1 1]
E5b-Q code (10230 bits, 1msec, 14-stage Gold code)		E5a-Q code (10230 bits, 1msec, 14-stage Gold code)	
Polynomial_1	X ¹⁴ +X ¹³ +X ¹¹ +X ⁴ +1	Polynomial_1	X ¹⁴ +X ⁸ +X ⁶ +X+1
Initial State_1	[1 1 1 1 1 1 1 1 1 1 1 1 1]	Initial State_1	[1 1 1 1 1 1 1 1 1 1 1 1 1]
Polynomial_2	X ¹⁴ +X ¹² +X ⁹ +X ⁸ +X ⁵ +X ² +1	Polynomial_2	X ¹⁴ +X ¹² +X ⁸ +X ⁷ +X ⁵ +X ⁴ +1
Initial State_2	[1 1 0 0 0 0 0 0 0 0 0 1 0 0]	Initial State_2	[0 1 1 0 1 1 0 0 1 0 1 0 1 0]

- All four codes are linear codes, can be generated by a 28 order polynomial
- The obtained code generator polynomials can be factorized
- Ambiguity of sign flipping is solved
- Gold codes!



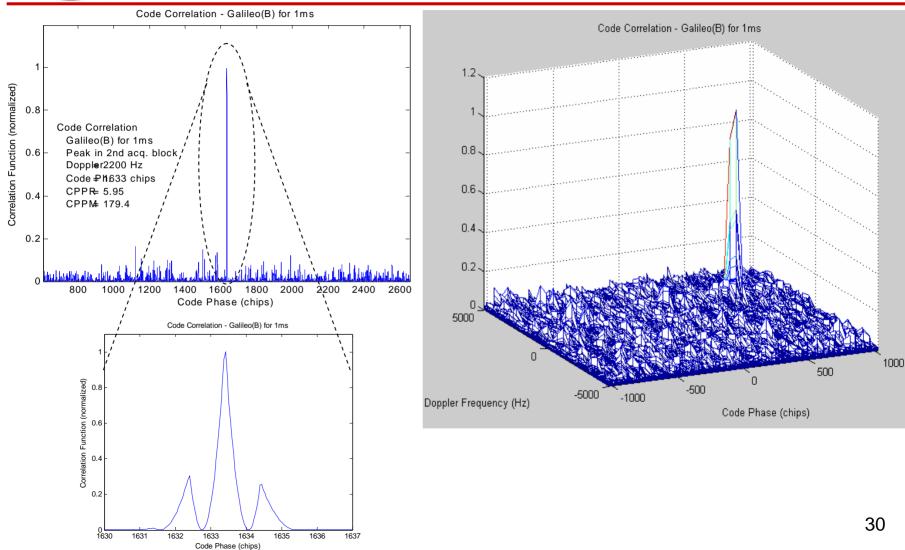
Applying the Codes for Acquisition and Tracking Real Broadcast signal



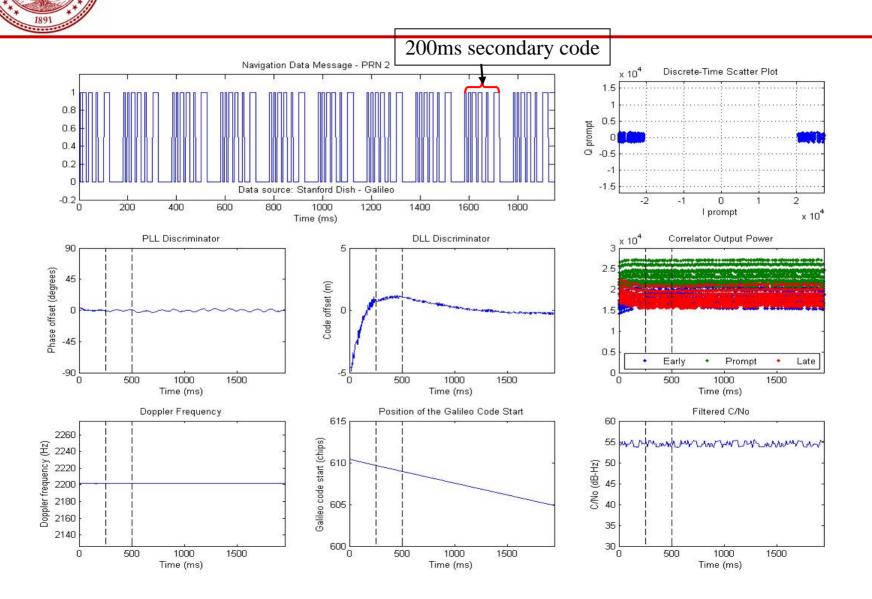
Software receiver block diagram



Acquisition of Broadcast L1-C Signal

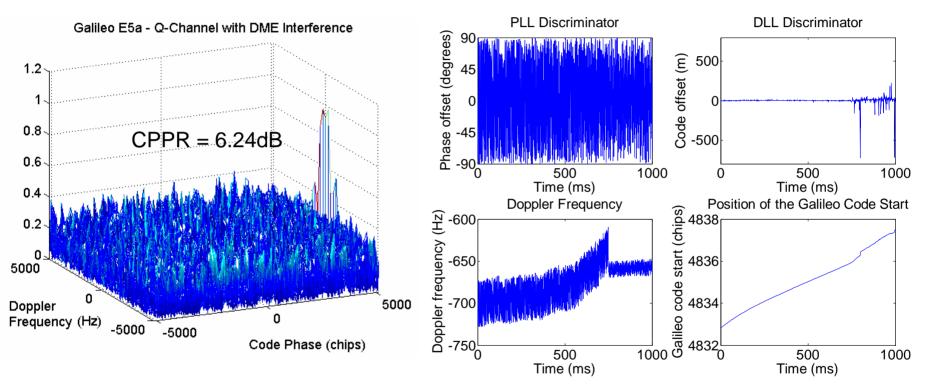


Tracking of Broadcast L1-C Signal





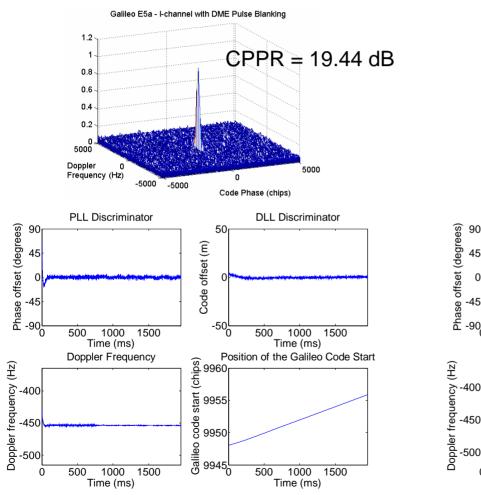
Acquire and Track Raw Data without any Signal Processing

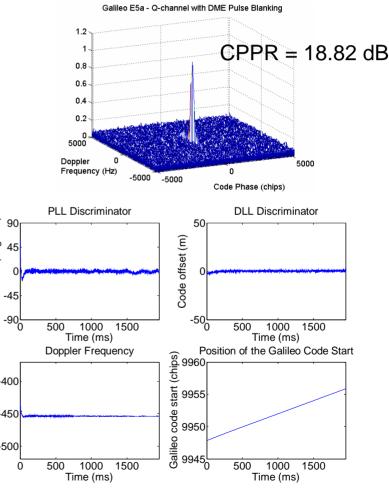


- Acquisition "successful" but tracking does not converge too much error in initial Doppler estimate
- Doppler should be -450 Hz, but PLL incorrectly locks onto -650 Hz; poor aiding to DLL causes periodic large discriminator errors



Acquire and Track Pulse-blanked Data

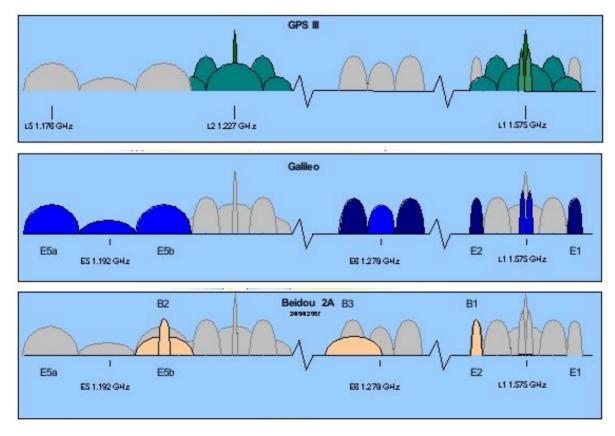






Compass Frequency Plan Compared with GPS and Galileo

- On April 13 China launched the first MEO satellite in its Compass GNSS system, 21,550 km above the Earth. (GPS: 20,184 km; Galileo: 23,222 km)
- It has been transmitting on L1, E6 and E5b.



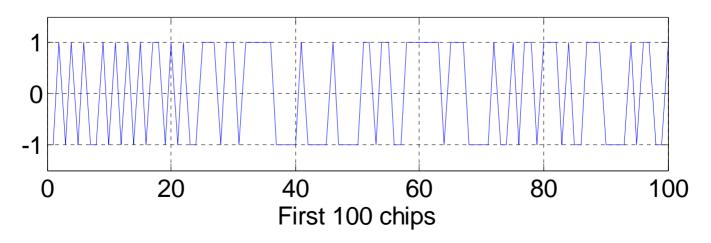
http://www.insidegnss.com/compass.php



Decoded Chinese Compass PRN Codes

The L1 PRN code is decoded to be an 11-order Gold Code.

L1 code (2046 bits, 1msec, 11-stage Gold code)			
Polynomial_1	$X^{11}+X^{10}+X^9+X^8+X^7+X+1$		
Initial State_1	[0101010101010]		
Polynomial_2	$X^{11}+X^9+X^8+X^5+X^4+X^3+X^2+X+1$		
Initial State_2	[0000001111]		





Summary

- Decode all broadcast codes in L1, E6, E5a and E5b frequency bands of GIOVE-A test satellite.
- Decode Compass L1 code.
- Prove codes to be truncated Gold Codes
- Find code generators linear shift feedback register
- Implement the code generators in software receivers
- Be able to acquire and track real broadcast signals
- Interference mitigation is desirable and beneficial for E5 signals



Questions?

